

Aldor: Interfaces

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Outline

- Inter-language communication
- Additional category concepts
- Object oriented programming
- Ex post facto extensions
- Specializing generic constructors

Interlanguage Communication

Call *to* and be called *by* programs in other languages/systems.

- Axiom
- C
- C++
- Fortran
- Maple

Importing a function from C

```
--  
-- arigato.as: A main AxiomXL program calling the C function 'nputs'.  
--  
#include "axllib.as"  
SI ==> SingleInteger;  
  
import { nputs: (SI, String) -> SI } from Foreign C;  
import from SI;  
  
nputs(3, "Arigato gozaimasu!");
```

Exporting a function to C

```
...
-- aside.as: An Aldor function made available to C.
--
#include "axllib.as"

SI ==> SingleInteger;

export { lcm: (SI, SI) -> SI } to Foreign C;

_lcm(n: SI, m: SI): SI == (n quo gcd(n,m)) * m;
```

```
/*
 * cside.c: A main C program calling the Aldor function 'lcm'.
 */
#include "foam_c.h"

extern FISInt lcm (FISInt, FISInt);

int
main()

    printf("The lcm of 6 and 4 is %d\n", lcm(6,4));
    return 0;
```

Category Defaults

- Often certain operations are equivalent to certain combination of others. Aldor therefore allows default values to be specified:

```
define AbelianGroup: Category == with {
  0: %;
  +: (% , %) -> %;
  -: % -> %;
  -: (% , %) -> %;
  *: (Integer, %) -> %;

  default {
    (a: %) - (b: %): % == a + (-b);
    (n: Integer) * (a: %): % ==
      power(a, n)$BinaryPowering(% , 0, +);
  }
}
```

- Defaults are late binding. A domain satisfying `AbelianGroup` will provide its own version of an export in preference to a default.

```

define OrderedSet: Category == Set with {
  <: (% , %) -> Boolean;
  >: (% , %) -> Boolean;
  <=: (% , %) -> Boolean;
  >=: (% , %) -> Boolean;
  default {
    (a: %) > (b: %): Boolean == b < a;
    (a: %) <= (b: %): Boolean == a < b or a = b;
    (a: %) >= (b: %): Boolean == a > b or a = b;
  }
}

...
default {
  (a: %) < (b: %): (Boolean, %) == (a < b, b);
  (t: Boolean, a: %) < (b: %): (Boolean, %) == ...
  ...
}

...
}

...: OrderedSet == add {
  s = t == ...
  t < t == ...
  -- Rest given by cat defaults
}

... f t0 < t1 <= t2 = t3 < t4 then ...

```


Conditional Categories

- Typically, additional assertions can be made when more properties are known to be satisfied.

Aldor therefore provides conditional construction of domains and categories.

```
UnivariatePolynomialCategory(R: Ring): Category == Module R with {
    monomial: (R, Integer) -> %;
}
if R has Field then EuclideanDomain
}
SparseUnivariatePolynomial(R: Ring): UnivariatePolynomialCategory(R) == {
    ...
    if R has Field then {
        gcd(p: %, q: %): % == ...
    }
}
```

```
Define DirectProductCategory(S: Set): Category == Set with {
  apply: (% , SingleInteger) -> S;
  if S has Semigroup then Semigroup;
  if S has Monoid then Monoid;
  if S has Group then Group;
  if S has Ring then Ring;
  if S has Finite then Finite;
  if S has Ordered then Ordered;
```

Categories as Views

MyDom(K: Field): Join(Ring, Algebra K, VectorSpace K) == add ...

FM(M: Monoid) == ...

FR(R: Ring) == ...

FA(K: Field, A: Algebra K) == ...

FM(MyDom Rational)

FR(MyDom Z7)

FA(IntegerMod 7, MyDom IntegerMod 7)

The Object Constructor

- From the Aldor library:

```
Object(C: Category): with {
  object:      (T: C, T) -> %;
  avail:      % -> (T: C, T);
}
== add {
  Rep == Record(T: C, val: T);
  import from Rep;

  object(T: C, t: T) : % == per [T, t];
  avail (ob: %) : (T: C, T) == explode rep ob;
}
```

- E.g.

```
import from String, Integer, List Integer;

boblis: List Object BasicType := [
    object (String,      "Ahem!"),
    object (Integer,     42),
    object (List Integer, [1,2,3,4])
];

bobfun(T: BasicType, t: T) : () ==
    print << "This prints itself as: " << t << newline;
for bob in boblis repeat bobfun(availl bob)
```

Post Facto Extensions

- If T is a type-valued constant, then its meaning may be extended with a definition of the form:

`extend T: C == D`

- C is a new category to which T will now belong
- D is a domain providing the newly required exports.

- Example: `Integer` can be made to satisfy the new categories `FancyOutput` and `DifferentialRing` by providing appropriate extensions:

```
extend Integer: FancyOutput == add {  
    box(n: %): BoundingBox == [1, ndigits n, 0, 0]  
}  
  
extend Integer: DifferentialRing == add {  
    differentiate(n: %): % == 0  
}
```

- This generalizes to the extension of functions which compute types.

Specializing Generic Constructors

```
SI==> SingleInteger;

QuadraticForm(n: Integer, R: Ring):

  Join(Algebra R, MatrixCategory R) with {
    apply: (% , Vector(n, R)) -> R;
  }

== SquareMatrix(n, R) add {
  Rep == SquareMatrix(n, R);
  import from Rep;

  apply(q: %, v: Vector(n, R)): R == {
    r: R := 0;
    for i in 1..n::SI repeat {
      s: R := 0;
      for j in 1..n::SI repeat s := s + rep(q)(i, j)*v(j);
      r := r + v(i) * s;
    }
    r;
  }
}
```



```

 CliffordAlgebra(n: Integer, K: Field, Q: QuadraticForm(n, K)):
    Join(Ring, Algebra K, VectorSpace K) with {
        e: Integer -> %;
        term: (K, List Integer) -> %;
        coef: (% , List Integer) -> K;
    }
}

== add {
    Rep == Array(K);

    import from Rep;
    import from SI;
    import from Vector(n, K);

    local nsi: SI == n::SI;
    local nel: SI == 2^nsi;
    local qee: Array K := [ Q e i for i in 1..n ];

    local new(): Rep == new(nel, 0);
    local apply(a: %, i: SI): K == rep(a).(i);

    dimension: Integer == 2^n;
}

```

```

0: % == per new();
1: % == { local z := new(); z.1 := 1; per z }

(x: %) + (y: %): % == per [ x.i + y.i for i in 1..nel ];
(x: %) - (y: %): % == per [ x.i - y.i for i in 1..nel ];
(k: K) * (x: %): % == per [ k * x.i for i in 1..nel ];
(x: %) * (k: K): % == per [ x.i * k for i in 1..nel ];
- (x: %): % == per [ -x.i for i in 1..nel ];
(x: %) / (k: K): % == per [ x.i / k for i in 1..nel ];

coerce(m: SI): % == { local z := new(); z.1 := m::K; per z }
coerce(m: Integer): % == { local z := new(); z.1 := m::K; per z }
coerce(k: K): % == { local z := new(); z.1 := k; per z }

(x: %) = (y: %): Boolean == {
  for i in 0..nel repeat if x.i ~= y.i then return false;
  true
}

```

```

(x: %) ^ (n: Integer): % == {
    n < 0 => error "Cannot take inverse";
    n = 0 => 1;
    n = 1 => x;
    n = 2 => x*x;
    odd? n => x*x^(n-1);
    x^(n quo 2)^2;
}

-- Basis elements are e1..en.
e(i: Integer): % == {
    i >= n => error "No such basis element";
    iz: SI := 2^i;
    local z := new(); z.iz :=1; per z
}

```

```

-- Returns coef and index after swapping e1 into correct order.
-- Indices k in the array are  $1..2^n$ .
-- ej => k-1 =  $2^{(j-1)}$ . Or together.

local canonIndex(e1: List Integer): (K, SI) == {
  -- 1. Check input
  for e in e1 repeat
    e >= n => error "No such basis element";

  -- 2. Apply identity ei ej = -ej ei, i ~ j.
  -- Bubble sort keeps the swaps obvious. OK because n is small.
  ea: Array SI := [ e::SI for e in e1 ];
  ne := #ea;

  exchanges: SI := 0;
  changed := true;
  while changed repeat {
    changed := false;
    for i in 1..ne-1 repeat
      if ea(i) > ea(i+1) then {
        t := ea(i); ea(i) := ea(i+1); ea(i+1) := t;
        exchanges := exchanges + 1;
        changed := true;
      }
    }
  }
  coef: K := if odd? exchanges then -1 else 1;
}

```

```

-- 3. Apply identity  $ei ei = Q(ei)$ , and compute index.
kminus1 := 0;
for e in ea repeat {
  jminus1 := e::SI - 1;
  if bit(kminus1, jminus1) then {
    coef := coef * Qee(jminus1+1);
    kminus1 := kminus1 - 2^jminus1;
  }
  else
    kminus1 := kminus1 + 2^jminus1;
}
(coef, kminus1 + 1)
}

```

```

local canonIndexProd(ix: SI, iy: SI): (K, SI) == {
  c: K := 1;
  izminus1: SI := ix - 1;
  izminus1: SI := iy - 1;
  for i in 0..nsi-1 | bit(izminus1, i) repeat {
    -- Apply rule  $e_i * e_j = -e_j * e_i$  for  $i \neq j$ 
    k: SI := 0;
    for j in i+1..nsi-1 | bit(izminus1, j) repeat k := k + 1;
    for j in 0..i-1 | bit(izminus1, j) repeat k := k + 1;
    if odd? k then c := -c;

    if bit(izminus1, i) then {
      c := c * Qee(i+1);
      izminus1 := izminus1 - 2^i
    }
    else
      izminus1 := izminus1 + 2^i;
  }
}
(c, izminus1 + 1)
}

```

```

(x: %)*(y: %): % == {
  local z := new();
  for ix in 1..nel | x.ix ~= 0 repeat
    for iy in 1..nel | y.ix ~= 0 repeat {
      (c, iz) := canonIndexProd(ix, iy);
      if c ~= 0 then
        z.iz := z.iz + x.ix * y.iy / c
      }
    }
  per z
}

term(c: K, e1: List Integer): % == {
  (ic, ix) := canonIndex(e1);
  local z := new(); z.ix := ic * c; per z
}

coef(z: %, e1: List Integer): K == {
  (ic, ix) := canonIndex(e1);
  ic = 0 => error "Asking for coefficient of []";
  z.ix/ic
}

```

```

local writeTerm(o: TextWriter, coef: K, ix: SI): () == {
  ix = 1 => o << coef;
  o << coef;
  o << " ";
  for i in 1..n::SI repeat {
    if bit(ix, i-1) then o << "e_" << i
  }
}

(o: TextWriter) << (x: %): TextWriter == {
  nout: SI := 0;
  for i in 1..nel repeat {
    c := x.i;
    if c ~= 0 then {
      if nout > 0 then o << " + ";
      nout := nout + 1;
      writeTerm(o, c, i);
    }
  }
  if nout = 0 then o << "0";
  o
}
}

```


Specializing CliffordAlgebra (Axiom)

```
-- The complex numbers as a Clifford Algebra
K := FRAC POLY INT
(1) Fraction Polynomial Integer
qf := QFORM(1, K) := quadraticForm(matrix([[ -1]]))$(SQMATRIX(1, K))
(2) [ - 1]
K := CLIF(1, K, qf)
(3) CliffordAlgebra(1, Fraction Polynomial Integer, MATRIX)
e := e(1)$C
(4) e
      1
i := a + b * i
(5) a + b e
      1
```

$$r := c + d * i$$

$$(6) \quad c + d e_1$$

$$: * y$$

$$(7) \quad - b d + a c + (a d + b c) e_1$$

-- The quaternions as a Clifford Algebra

qf:=QFORM(2, K) :=quadraticForm matrix([[-1, 0], [0, -1]])\$(SQMATRIX(2, K))

$$(8) \quad \begin{array}{ccc|c} & -1 & 0 & | \\ & | & & | \\ & | & & | \\ (8) & | & & | \\ & 0 & -1 & | \end{array}$$

:= CLIF(2, K, qf)

(9) CliffordAlgebra(2, Fraction Polynomial Integer, MATRIX)

:= e(1)\$H

$$(10) \quad \begin{array}{c} e \\ 1 \end{array}$$

:= e(2)\$H

$$(11) \quad \begin{array}{c} e \\ 2 \end{array}$$

:= i * j

$$(12) \quad \begin{array}{cc} e & e \\ 1 & 2 \end{array}$$

$:= a + b * i + c * j + d * k$

$$(13) \quad a + b e_1 + c e_2 + d e_1 e_2$$

$r := e + f * i + g * j + h * k$

$$(14) \quad e + f e_1 + g e_2 + h e_1 e_2$$

$s * y$

(15)

$$- d h - c g - b f + a e + (c h - d g + a f + b e) e_1$$

+

$$(- b h + a g + d f + c e) e_2 + (a h + b g - c f + d e) e_1 e_2$$

-- The exterior algebra on a 3 space.

```
qf := QFORM(3, K) := quadraticForm(0::SQMATRIX(3,K))
```

```
      |0  0  0|
      |  1  1|
(18)  |0  0  0|
      |  1  1|
      |0  0  0|
```

```
Ext := CLIF(3,K,qf)
```

```
(19)  CliffordAlgebra(3,Fraction Polynomial Integer, MATRIX)
```

```
      := e(1)$Ext
```

```
(20)  e
      1
```

```
      := e(2)$Ext
```

```
(21)  e
      2
```

:= e(3)\$Ext

$$(22) \quad e_3$$

:= x1*i + x2*j + x3*k

$$(23) \quad x_1 e_1 + x_2 e_2 + x_3 e_3$$

:= y1*i + y2*j + y3*k

$$(24) \quad y_1 e_1 + y_2 e_2 + y_3 e_3$$

+ y

$$(25) \quad (y_1 + x_1)e_1 + (y_2 + x_2)e_2 + (y_3 + x_3)e_3$$

* y + y * x

$$(26) \quad 0$$

-- In n space, a grade p form has a dual n-p form.

-- In particular, in 3 space the dual of a grade 2 element identifies
e1*e2->e3, e2*e3->e1, e3*e1->e2.

```
dual2 a ==  
  coefficient(a, [2,3])$Ext * i +  
  coefficient(a, [3,1])$Ext * j +  
  coefficient(a, [1,2])$Ext * k
```

-- The vector cross product is then given by
dual2(x*y)

$$(28) \quad (x_2 y_3 - x_3 y_2)e_1 + (-x_1 y_3 + x_3 y_1)e_2 + (x_1 y_2 - x_2 y_1)e_3$$

-- The Dirac Algebra

:= FRAC INT

(29) Fraction Integer

S: SPMATRIX(4, K) := [[1, 0, 0, 0], [0, -1, 0, 0], [0, 0, -1, 0], [0, 0, 0, -1]]

```
| 1  0  0  0 |
|   |   |   |
| 0 -1  0  0 |
|   |   |   |
| 0  0 -1  0 |
|   |   |   |
| 0  0  0 -1 |
```

(30)

f: QFORM(4, K) := quadraticForm g

:= CLIF(4, K, qf)

(32) CliffordAlgebra(4, Fraction Integer, MATRIX)

-- The usual notation is $\gamma_{sup i}$.
 $\gamma_{am} := [e(i)]_{D}$ for i in $1..4]$

$$(33) \quad [e, e, e, e] \\ 1 \quad 2 \quad 3 \quad 4$$

-- There are various contraction identities of the form
 -- $g(1,t)*\gamma_{am}(1)*\gamma_{am}(m)*\gamma_{am}(n)*\gamma_{am}(r)*\gamma_{am}(s)*\gamma_{am}(t) =$
 -- $2*(\gamma_{am}(s)\gamma_{am}(m)\gamma_{am}(n)\gamma_{am}(r) + \gamma_{am}(r)\gamma_{am}(n)*\gamma_{am}(m)*\gamma_{am}(s))$
 -- where the sum over l and t is implied.

-- Verify this identity for $m=1, n=2, r=3, s=4$
 $n := 1; n := 2; r := 3; s := 4;$

$hs := \text{reduce}(+, [\text{reduce}(+, _$
 $[g(1,t)*\gamma_{am}(1)*\gamma_{am}(m)*\gamma_{am}(n)*\gamma_{am}(r)*\gamma_{am}(s)*\gamma_{am}(t) -$
 $\text{for } l \text{ in } 1..4]) \text{ for } t \text{ in } 1..4])$

$$(35) \quad - 4e e e e \\ 1 \quad 2 \quad 3 \quad 4$$

$hs := 2*(\gamma_{am} s * \gamma_{am} m * \gamma_{am} n * \gamma_{am} r + \gamma_{am} r * \gamma_{am} n * \gamma_{am} m * \gamma_{am} s)$

$$(36) \quad - 4e e e e \\ 1 \quad 2 \quad 3 \quad 4$$